Introduction: The main purpose of this project is to implement a floating point calculator that can do the following functions:

- Multiply two 2 bits number and display result when "=" button is pressed
- Divide two 32 bit numbers and display result when "=" button is pressed
- Change the sign of the number and display the result when the "=" sign is pressed

The FP calculator is implemented using a Keypad and an Altera DE2 FPGA development board. The results are being displayed on a 7 segment display. The results are displayed in Hex. The keypad is attached to the DE2 board using the GPIO pins. The Keypad is restricted to 20 Keys. The design is made using Quartus and everything is programmed in AHDL.

The code for the design can be found under appendix at the end of the report.


Figure 1: Here is the set up for the project


Figure 2 : The keypad that consist of 20 Keys


Figure 3: The seven segment display the Hex of the number being pressed. In this case 2 is being pressed, so it shows $0 \times 4000$

For this project the IEEE 754, 32 bit floating point standard is being used. The numbers being displayed on the 7 - Segment display are Hexadecimal number as it can be seen from figure 3.

Round Even and Round Odd standard: When the division leaves the reminder and the result is not exact we need to decide how we are going to round the result. If the reminder is 0 , it means that the number is exact. But if the reminder is in between $1-3$ the result is truncated to nearest integer value. If the reminder is in between 5-9, the result is incremented by adding 1 to it. The diagram below shows how it works.

## 000 : Result is Exact



100


For the case when the result is exactly in between the truncated and incremented values we need to choose what we would like to make the result.

If the result is chosen to be even then it is called even round. For example for the number like 378.5 if the even round about is applied the result will be 386, which is even. For number like 377.5 if the even round about is applied the result will end up being value incremented by 1 , which is 378 . This same thing can be applied to odd round about. In that case 377.5 will be 377 .

So with even or odd round about half the time result will end up being truncated to the nearest even /odd value and other half time the result will be incremented by 1 to the nearest even/odd value.

## Data Flow:

The Quartus Floating point calculator project is consist of 4 different modules that falls under the one top level design.

Here is the hierarchy.


Figure 4 : Hierarchy for the FP calculator design
The four modules are included in the top level design as include files. The design runs on 27 MHz clock.
Here are the outputs and inputs to the design.
Inputs: 27 MHz clock, row bits from the Keypad
Output: Column bits to the Keypad, Seven segment display
The four modules inputs and outputs are connected to the FPCalc.tdf via wires. All four modules are instantiated in the main design.

Calculator consists of two 32 bit floating point registers. Xx and yy . The final result for any operation is always stored in xx register.

Multiplication unit: The instantiation used for the FPMult in the design is MultUnit. The $x x$ and $y y$ registers of MultUnit are wired to the $x x$ and $y y$ register in the design respectively. The last key in the first row on the keypad is associated with the multiplication operation. "op" is a filp flop , that can assume only two values, either 0 or 1 . The value 0 is assigned to the multiplication operation in this design.

Division Unit: The instantiation used for the FPdivide in the design is DivUnit. The xx and yy registers of DivUnit are wired to the $x x$ and yy register in the design respectively. The last key in the second row on the keypad is associated with the division operation. The value 1 is assigned to the division operation for the "op" FF in this design.

Sign Change: This is the simplest of all operation. It needs very few lines of code. That is why is designed directly in the top level design. The last key in the thirst row is associated with the sign change operation.

The last key in the fourth row is equal to operation.

Adding new operations: if the new operation like addition and subtraction are to be added to the calculator design we will need two new modules. FPSub and FRAdd. Two separate keys will be needed on the calculator that represents Addition and Subtraction.

The design of these two modules will be quite similar to the Multiplication or division unit only the state machine will be simpler for this operation. We will need two floating point registers that we can use for the add / sub operations and that will be wired to the top level design the same way we wired the other two modules.

Control Units: Here are the state diagrams for the controller, multiplication, division and Keypad.


Figure 5 : State Diagram for Controller


Figure 6 : State Diagram for multiplication


Figure 7 : State Diagram for Division


Figure 8: State Diagram for Keypad

## Summary:

The 32 bit floating point calculator is capable to do $x, \div, \pm$, and $=$ operations. It uses a simple 3 state controller to implement this and displayed the results to the seven segment display on the board.

Here are the few things that can improve the current design:

- Add more function like Addition and Subtraction
- Add reset to the design instead of defaulting the registers values to 1
- Add an ON/OFF switch
- Being able to store the result after doing one operation to do more operations on it

Other than these we can add in a lot more other features to the design to make it better and more versatile.

## Appendix:

## FPclac.tdf

```
include "FPmult.inc";
include "FPdivide.inc";
include "Hto7seg.inc";
include "Keypad.inc";
Subdesign FPCalc
    (
    Clk27M :Input;
        r[0..3] :Input;
        c[0..3]
        :Output;
    SevenSeg[55..0] :Output; -- results showm on 7-seg display
    )
Variable
    xx[31..0], yy[31..0] :Dff;
    dig[7..0] :Hto7seg;
            op
:Dff;
    MultUnit :FPmult;
    DivUnit :FPdivide;
    Keyinterface :Keypad;
        key1,key2,key3,key4 :Node;
            keyCS,keyDiv,keyMult :Node;
            keyEq,keyUnused,newkey :Node;
    State :Machine
        With States(Init,Keys,Oper);
Begin
    xx[].clk = Clk27M;
    yy[].clk = Clk27M;
    State.clk = Clk27M;
        op.clk = Clk27M;
    MultUnit.Clk = Clk27M;
    MultUnit.xx[] = xx[];
    MultUnit.yy[] = yy[];
    DivUnit.Clk = Clk27M;
    DivUnit.xx[] = xx[];
    DivUnit.yy[] = yy[];
        Keyinterface.Clk27 = Clk27M;
        Keyinterface.r[] = r[];
        newkey = Keyinterface.nk;
        c[] = Keyinterface.c[];
-- prod[].d = MultUnit.prod[];
-- quo[].d = DivUnit.quo[];
-- prod[].clk = CLK27M;
-- quo[].clk = CLK27M;
-- MultDone = MultUnit.MultDone;
```

$$
\begin{aligned}
& \text { dig7.in[] = xx[31..28]; } \\
& \text { dig6.in[] = xx[27..24]; } \\
& \text { dig5.in[] }=x x[23 . .20] ; \\
& \text { dig4.in[] }=x x[19 . .16] ; \\
& \text { dig3.in[] = xx[15..12]; } \\
& \text { dig2.in[] }=x x[11 . .8] ; \\
& \text { dig1.in[] = xx[7..4]; } \\
& \text { dig0.in[] = xx[3..0]; } \\
& \text { SevenSeg[55..49] = dig7.out[]; } \\
& \text { SevenSeg[48..42] = dig6.out[]; } \\
& \text { SevenSeg[41..35] = dig5.out[]; } \\
& \text { SevenSeg[34..28] = dig4.out[]; } \\
& \text { SevenSeg[27..21] = dig3.out[]; } \\
& \text { SevenSeg[20..14] = dig2.out[]; } \\
& \text { SevenSeg[13..7] = dig1.out[]; } \\
& \text { SevenSeg[6..0] = dig0.out[]; }
\end{aligned}
$$

-- Decoded Keys for the Keypad
key1 $=$ Keyinterface. Rn[]$==0$ \& Keyinterface.Cn[]==0;
key2 $=$ Keyinterface.Rn[]==0 \& Keyinterface.Cn[]==1;
key3 = Keyinterface.Rn[]==0 \& Keyinterface.Cn[]==2;
key4 = Keyinterface.Rn[]==0 \& Keyinterface.Cn[]==3;
keyMult = Keyinterface.Rn[]==0 \& Keyinterface.Cn[]==4;
keyDiv = Keyinterface.Rn[]==1 \& Keyinterface.Cn[]==4;
keyCS = Keyinterface. $\mathrm{Rn}[\mathrm{]}==2$ \& Keyinterface. $\mathrm{Cn}[\mathrm{]}==4$; -- +/- (change-sign key)
keyEq = Keyinterface. $\mathrm{Rn}[\mathrm{]}==3$ \& Keyinterface. Cn[]$==4 ; \quad$-- $\quad$ (the equal key)
keyUnused = Keyinterface. Rn[]$!=0$ \& Keyinterface. Cn[]$!=4 ;-$ assuming bottom left $4 \times 3$ key area not used
-- The following is the main state machine -- most activities happen in the Keys state.
-- A value must be assigned to $x x[], y y[], ~ o p ~ a n d ~ n e x t ~ S t a t e ~ f o r ~ e v e r y ~ p o s s i b l e ~$
-- condition in each state WITHOUT OVERLAP (setting it to different things in
-- conditions that overlap so both are true). I think I have done that below,
-- but I have not reread it or tested it in any way.

## Case State Is

When Init =>
$x x[] . d=H " 3 F 800000 " ;-$ Initial value is 1.0
yy[].d = H"3F800000"; -- Initial value is 1.0
-- op will become 0 which specifies multiply
State = Keys;

When Keys =>
-- first handle the no keypress condition and unused keys (lower left 4x3)

If !newkey \# newkey \& keyUnused Then
$x x[] . d=x x[] ; y y[] . d=y y[] ;$ op.d = op; State = Keys;
End If;
-- the following section handles the multiply and divide keys
-- If newkey \& (keyMult \# keyDiv) Then
-- $\quad y y[] . d=x x[] ; x x[] . d=x x[] ;$ State $=$ Keys;
-- op.d = keyDiv; -- this sets op = 0 for Mult and op $=1$ for Div
-- End If;
If newkey \& keyMult Then

$$
y y[] . d=x x[] ; x x[] . d=x x[] ; \text { State }=\text { Keys; }
$$

op.d = B"0"; End If;
If newkey \& keyDiv Then
yy[]$. \mathrm{d}=\mathrm{xx}[] ; \mathrm{xx}[] . \mathrm{d}=\mathrm{xx}[]$; State $=$ Keys;
op.d = B"1"; End If;
-- Handle the change-sign keypress
If newkey \& keyCS Then yy[].d = yy[]; op.d = op; State = Keys; xx[].d = (!xx31,xx[30..0]);
End If;
-- Handle the number keys
If newkey \& key1 Then yy[].d = yy[]; op.d = op; State = Keys; xx[].d = H"3F800000"; End If; -- 1.0
If newkey \& key2 Then yy[].d = yy[]; op.d = op; State = Keys; xx[].d = H"40000000"; End If; -- 2.0
If newkey \& key3 Then yy[].d = yy[]; op.d = op; State = Keys; xx[].d = H"40400000"; End If; -- 3.0
If newkey \& key4 Then

$$
\text { yy[].d = yy[]; op.d = op; State = Keys; xx[].d = H"41200000"; End If; -- } 10.0
$$

-- next handle the equal sign
If newkey \& keyEq Then
yy[].d = yy[]; op.d = op; State = Oper;
xx[]$. \mathrm{d}=\mathrm{xx}[]$; MultUnit.start $=$ !op; DivUnit.start $=\mathrm{op}$;
End If;
When Oper =>
yy[].d = yy[]; op.d =op;
If (MultUnit.MultDone \# DivUnit.DivDone) Then -- only the operation started can
assert done
State $=$ Keys; $x x[] . d=$ MultUnit.prod[] + DivUnit.quo[]; -- one will be zero so not harm
other
Else
State $=$ Oper; $\mathrm{xx[].d=xx[];}$
End If;

End Case;
End;

## FPmult.tdf

```
Subdesign FPmult
    l
    xx[31..0], yy[31..0], Clk, Start :Input; -- fp numbers to multiply & start signal
    prod[31..0], MultDone :Output; -- product and a done signal
    )
Variable
    S, E[9..0], A[23..0], B[23..0] :Dff;
    R, SB, P[24..0], MultDone :Dff;
    Done, Rneed :Node;
    State :Machine
        With States(Wait,Step,Shift);
Begin
    S.clk = Clk; E[].clk = Clk; A[].clk = Clk; B[].clk = Clk; State.clk = Clk;
    R.clk = Clk; SB.clk = Clk; P[].clk = Clk; MultDone.clk = Clk;
    Done = (A[]==0)&!P24; Rneed = R&SB # P0&R&!SB;
    prod[] = (S,E[7..0],P[22..0]);
    Case State Is
        When Wait =>
            If Start Then
                    S.d= xx[31] $ yy[31]; A[].d = (1,xx[22..0]); B[].d = (1,yy[22..0]);
                    Else
                    State = Wait; --Note: All Dffs go to zero so prod[] becomes zero
                    End If;
                When Step =>
                    S.d = S; R.d = PO; SB.d = SB # R;
                    A[].d = (0,A[23..1]); B[].d = B[];
                        If !Done Then
                    P[].d = (0,P[24..1]) + (0,A0&B[]); E[].d = E[] + 1;
                            State = Step;
                    Else
                    E[].d = E[];
                        If !Rneed Then
                                P[].d = P[]; MultDone.d = B"1"; State = Wait;
                            Else
                                    P[].d = P[] + 1; State = Shift;
                                    End If;
                            End If;
                When Shift =>
                        S.d = S; MultDone.d = B"1"; State = Wait;
                If P24 Then
```


## Floating Point Calculator

```
\(E[] . d=E[]+1 ; P[] . d=(0, P[24 . .1]) ;\)
Else
    E[].d = E[]; P[].d = P[];
End If;
```

    End Case;
    End;

## FPdivide.tdf

```
Subdesign FPdivide
    (
    xx[31..0], yy[31..0], Clk, Start :Input; -- fp divisor, dividend & start signal
    quo[31..0], DivDone :Output; -- quotient and a done signal
    )
Variable
    S, E[9..0], A[24..0], B[24..0] :Dff;
    Q[23..0], DivDone :Dff;
    Neg, Done, Rneed, AmB[24..0] :Node;
    State
                                :Machine
        With States(Wait,Step);
Begin
    S.clk = Clk; E[].clk = Clk; A[].clk = Clk; B[].clk = Clk;
    Q[].clk = Clk; DivDone.clk = Clk; State.clk = Clk;
    Done = Q23; Rneed = A[]>B[] # A[]==B[]&Q0;
    AmB[] = A[] - B[]; Neg = AmB24; quo[] = (S,E[7..0],Q[22..0]);
    Case State Is
        When Wait =>
                            If Start Then
                                    S.d = xx[31] $ yy[31]; A[].d = (01,yy[22..0]); B[].d = (01,xx[22..0]);
                                    E[].d = (0,yy[30..23]) - (0,xx[30..23]) + 127 + 24; State = Step;
                                    Else
                                    State = Wait; --Note: All Dffs go to zero so prod[] becomes zero
                                    End If;
                When Step =>
                            S.d = S; B[].d = B[];
            If !Neg Then A[].d = (AmB[23..0],0); Else A[].d = (A[23..0],0); End If;
                    If !Done Then
                        E[].d = E[]-1; Q[].d = (Q[22..0],!Neg); State = Step;
                            Else
                            E[].d = E[]; Q[].d = (Q[] + (0,Rneed)); DivDone.d = B"1"; State = Wait;
                            End If;
    End Case;
End;
Hto7seg.tdf
```

```
Subdesign Hto7seg
    (
    in[3..0] :Input; \%hex value in\%
    out[6..0] :Output; \%7 segents out\%
    )
Variable
    zero,one,two,three,four,five,six,seven,
    eight,nine, AA, BB,CC,DD,EE,FF :node;
Begin
    zero \(=\) in[]==0;
    one \(=\) in [] \(==1\);
    two \(=\) in[] \(==2\);
    three \(=\) in []\(==3\);
    four \(=\) in []\(==4\);
    five \(=\) in []\(==5\);
    six = in[]==6;
    seven \(=\) in []\(==7\);
    eight \(=\) in []\(==8\);
    nine \(=\) in []\(==9\);
    \(A A=\operatorname{in}[]==10\);
    \(B B=\operatorname{in}[]==11 ;\)
    CC = in[]==12;
    DD = in[]==13;
    \(\mathrm{EE}=\mathrm{in}[]==14\);
    \(\mathrm{FF}=\mathrm{in}[]==15\);
```

    out0 = !(zero \# two \# three \# five \# six \# seven \# eight \# nine \# AA \# CC \# EE \# FF);
    out1 = !(zero \# one \# two \# three \# four \# seven \# eight \# nine \# AA \# DD);
    out2 = !(zero \# one \# three \# four \# five \# six \# seven \# eight \# nine \# AA \# BB \# DD);
    out3 = !(zero \# two \# three \# five \# six \# eight \# BB \# CC \# DD \# EE);
    out4 = !(zero \# two \# six \# eight \# AA \# BB \# CC \# DD \# EE \# FF);
    out5 = !(zero \# four \# five \# six \# eight \# nine \# AA \# BB \# CC \# EE \# FF);
    out6 = !(two \# three \# four \# five \# six \# eight \# nine \# AA \# BB \# DD \# EE \# FF);
    End;

## Keypad.tdf

Subdesign Keypad
1
Clk27 : Input;
$r[0 . .3] \quad$ : Input ;
--ro[0..3],st0o : Output;
c[0..3] : Output;
$\operatorname{Rn}[1.0]$ : Output;
Cn[2..0] : Output;
nk
: Output ;
)

Variable
\%s[2..0] : Dff;\%
ctr[19..0] : Dff;
nk,Rn[1..0],Cn[2..0] : Dff;
done
count
: Node ;
: Node ;
: Node;

State : Machine With
States(st0,st1,st2,st3,st4,st5,st6,st7);

Begin
--ro[] = r[];
--st0o = st0;

Ctr[].clk = Clk27;
$\operatorname{Ctr}[] . d=$ count \& (Ctr[].q+1);
Rn[].clk = Clk27;
Cn[].clk = Clk27;
nk.clk = Clk27;

Rn1.d = st1 \& (! r2 \# !r3) \# ! st1 \& Rn1;
Rn0.d = st1 \& (! r1 \# !r3) \# !st1 \& Rn0;

Cn[].d = Cn[].q \& (st6 \# st7 \# st0 );
--Cn0 = c1 \# c3 ;
--Cn1 = c2 \# c3 ;
--Cn2 = Cn2 \& !Cn2;
--
$--c 0=$ State $==s t 2 \& n k$;
--c1 = State $==s t 3 \& n k ;$
--c2 = State == st4 \& nk;
--c3 = State == st5;
\%c4 = c4 \& !c4;\%
done $=$ Ctr19;
\%r[0..3] = !(!c0 \& r[0..3] \& c0 \# !c1 \& r[0..3] \& c1 \# !c2 \& r[0..3] \& c2 \# !c3 \& r[0..3] \& c3
\# !c4 \& r[0..3]\& c4) ;\%
allones $=r 0$ \& r1 \& r2 \& r3;
nk.d = allones \& (st2 \# st3 \# st4) \# st5;
\%s2.d = st2 \& allones \& nk \# st3 \& allones \& nk \# st4 \& allones \& nk \# st3 \& !allones \# st4 \& !allones \# st5 \# st6 \& !allones \# st6 \& allones \# st7 \& !done ;
s1.d = st1 \& done \# st2 \& allones \& nk \# st2 \& !allones \# st3 \& allones \& nk \# st4 \& allones \& nk \# st5 \# st6 \& !allones \# st6 \& allones \# st7 \& !done ;
s0.d = st0 \& !allones \# st1 \& !done \# st2 \& !allones \# st4 \& !allones \# st6 \& allones \# st7
\& !done;\%

State.clk = Clk27;
case State Is
When st0 =>
If !allones then State = st1;
Else State = st0;
End If;
When st1 =>
count = b"1" ;
If done then State $=$ st2;
Else State = st1;
End If;
When st2 =>
c0 = b"1";
If allones then State = st6; Cn[].d=0; End If;
If !allones then State = st3; End If;
When st3 =>
c1 = b"1" ;
If allones then State = st6; Cn[].d=1; End If;
If !allones then State = st4; End If;
When st4 =>
c2 = b"1" ;
If allones then State = st6; Cn[].d=2; End If;
If !allones then State = st5; End If;
When st5 =>
c3 = b"1" ;
State = st6; Cn[].d = 3 \& allones \# 4 \& !allones;
When st6 =>
If allones then State = st7;
Else State = st6;
End If;
When st7 =>
count = b"1";
If done then State = st0;
Else State = st7;
End If;
End Case;
End;

